

Graph curves

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Plan:

I. Define and review graph curves;
present new results;
mention open questions

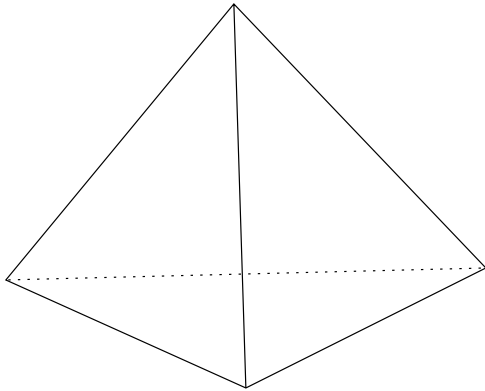
II. Use in GIT

representations of automorphism groups of curves

Definition 1 A graph curve C is a connected, reduced, projective algebraic curve whose only singularities are nodes, and such that the normalization of every irreducible component is \mathbf{P}^1 .

The graph is the *dual graph* of C :

- the irreducible components of C are the vertices of G
- the nodes of C are edges of G



Main reference: Bayer and Eisenbud. “Graph curves.” *Adv. Math.* **86** (1991) 1–30.

Running hypothesis: I will assume for the rest of this talk that G has no returning edges (because arithmetic genus 1).

Preliminaries: The genus of a graph curve is the Euler characteristic of the graph.

$$g(C) = \#edges - \#vertices + 1$$

In fact, once the edges are oriented, there is an isomorphism

$$H^0(C, \omega) \cong H^1(G, \mathbb{C})$$

$s \mapsto$ cochain whose value on edge
 from vertex α to vertex β
 representing node P is $\text{Res}_{C_{\alpha, P}} s$

[BE] study the multiplication maps

$$\begin{array}{ccc} H^0(\omega) & \otimes & H^0(\omega) & \rightarrow & H^0(\omega^2) \\ & & & & \downarrow \varphi \\ H^1(G, \mathbb{C}) & \otimes & H^1(G, \mathbb{C}) & \rightarrow & \text{Cochains}^1(G, \mathbb{C}) \end{array}$$

and show that for trivalent graphs, φ is an isomorphism, so

$$H^0(\omega^2) \cong \text{Cochains}^1(G, \mathbb{C}).$$

More generally, for trivalent graphs G one has

$$H^0(\omega^n) \cong \text{Coch}^1(G, \mathbb{C}) \oplus \bigoplus^{n-2} \text{Coch}^0(G, \mathbb{C})$$

Remark: in the stratification of $\partial\overline{\mathcal{M}}_g$ by number of nodes, trivalent graphs correspond to the 0-dimensional strata, i.e. they are as deep in the boundary as possible. Moreover, they do not deform/have no moduli.

Algorithm for canonical equations of some graph curve:

Let G be trivalent, 2-edge-connected, with no multiple edges (or returning edges). Write

$$S := \text{Sym}^\bullet H^0(\omega) \cong \text{Sym}^\bullet H^1(G, \mathbb{C})$$

$$S \hookrightarrow T := \text{Sym}^\bullet \text{Cochains}^1(G, \mathbb{C})$$

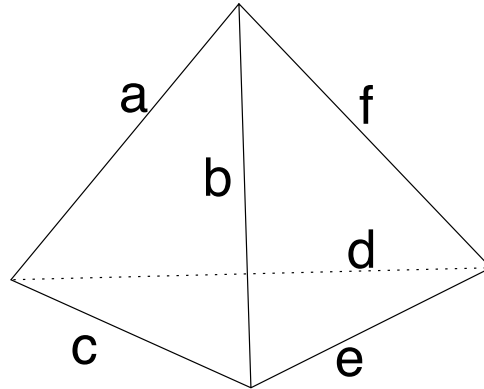
$$R = \bigoplus_n H^0(\omega^n)$$

$$\text{Then } \ker(S \rightarrow R) = I(C)$$

[BE]: I can be computed as the intersection of S with the ideal J of T generated by all monomials

ab , a, b disjoint edges

abc , a, b, c form a triangle



Example: tetrahedron

Disjoint edge pairs: ae , bd , cf

Triangles: abc , bef , cde , adf .

tetrahedron (continued)

$$J = (ae, bd, cf, abc, bef, cde, adf) \subset \mathbb{C}[a..f]$$

$S = \mathbb{C}[X, Y, Z]$ where

$$X := a + b + c$$

$$Y := b + e + f$$

$$Z := c + d + e.$$

$$\begin{aligned} \Rightarrow I &= \ker(S \rightarrow \mathbb{C}[a..f]/J) \\ &= (XYZ(X + Y + Z)). \end{aligned}$$

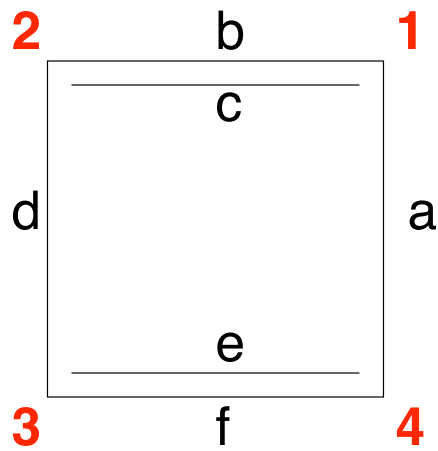
Algorithm [Swinarski]: 2-canonical equations of trivalent graph curves

Label each edge of the graph. If edges a, b, c meet at vertex α , then the ideal $I(C_\alpha)$ of the component C_α represented by α is given by

$$(ab + bc + ac) + (\{\text{edges not meeting vertex } \alpha\}).$$

Then

$$I(C) = \bigcap_{\text{vertices } \alpha} I(C_\alpha)$$



$$I(C_1) = (ab + bc + ac, d, e, f)$$

$$I(C_2) = (bc + cd + bd, a, e, f)$$

⋮

Check: genus of C is 3, $\deg \omega^2 = 8$ and Hilbert polynomial of I is $p(t) = 8t - 2$.

4 smooth components, 6 singular points:

$\{ [1:0:0:0:0:0], \dots [0:0:0:0:0:1] \}$

Generalize to n -canonical embeddings

- I have an algorithm for $n = 3$, but it's not in final form (the equations aren't obviously symmetric)
- the same strategy works for any $n \geq 2$, with the same beautification problem

Note: in contrast with the [BE] algorithm, the strategy for n -canonical equations works for any trivalent graph. Multiple edges are allowed, and there are no connectivity hypotheses on the graph

Also we can work one component at a time, and it is not necessary to choose a basis of $H^1(G, \mathbb{C})$.

Attempted application in GIT

Bayer–Morrison: Hilbert stability of an ideal I with respect to a 1-PS can be reformulated as whether a certain realization of the state polytope of I contains a specific point.

Thanks to Gfan, it is possible to compute state polytopes in small cases (e.g. 2-canonical embeddings of genus 3 curves). But this only tells us stability with respect to 1-PS inside one maximal torus.

Goal: find a way to conclude stability with respect to the full group \mathcal{G} after a small number of calculations.

Kempf: when y is unstable with respect to \mathcal{G} , there is a worst destabilizing 1-PS λ , and λ is preserved by a nontrivial parabolic $\mathcal{P} \subsetneq \mathcal{G}$. So if the stabilizer of y in \mathcal{G} is not contained in any proper parabolic subgroup of \mathcal{G} , then y is stable.

Kempf applied this to examples where the representation of the automorphism group is irreducible.

Morrison: if we at least knew that it decomposed into irreducibles all of multiplicity 0 or 1, then we would still “know” \mathcal{P} . A little linear programming may allow us to conclude that there is no destabilizing 1-PS.

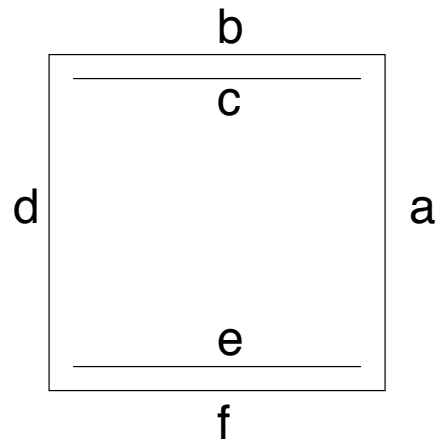
Definition 2 Say a (connected, reduced, projective) curve C with genus $g > 2$ is a 0,1-curve for $\text{Sym}^2 H^0(\omega^2)$ if in the decomposition of $\text{Sym}^2 H^0(\omega^2)$ into irreducibles with respect to the action of the automorphism group $\text{Aut}(C)$ of C , every irreducible has multiplicity 0 or 1.

Problem: Find 0,1-curves.

Theorem 3 (Swinarski) *There are no smooth complex 0,1-curves for $\text{Sym}^2 H^0(\omega^2)$.*

But nodal curves can have bigger and more abelian automorphism groups than smooth curves. Maybe we can find examples among trivalent graph curves?

An example will immediately illuminate a necessary condition:



Observe: $a + d$ and $b + c + e + f$ generate two copies of the trivial representation.

Proposition 4 *If there are two disjoint orbits of edges under $\text{Aut}(G)$, then $C(G)$ cannot be a $0,1$ -curve.*

Marston Conder has checked: no trivalent graphs on up to 768 vertices are 0,1-graph curves.

THE END
THANK YOU!

Example: the graph with two vertices, three edges.

$$I(C_x) = (y, \\ ac - ax - 2cx - 2x^2, \\ bc + bx + 2cx - 2x^2, \\ ab + ac + bc - ax + bx)$$

$$I(C_y) = (x, \\ ac - ay - 2cy - 2y^2, \\ bc + by + 2cy - 2y^2, \\ ab + ac + bc - ay + by)$$

$$I(C) = I(C_x) \cap I(C_y)$$