

Rob Kirby

1961 14 Aug - 8 Sept "Topology of Manifolds"
(Little red book)
Edited M.K. Fort Jr.

Decomposition spaces

n-Manifolds - M. Brown ∂M is collared
Stallings $F \rightarrow M^2$

Knots Fox - Quick-trip.

Poincaré Conjecture Zeeman + Stallings
Periodic Maps. for $n \geq 5$

1967 1 week

1968 1 week

1969 11-22 August "Topology of Manifolds"
Ed. Cantrell & C.H. Edwards

Triangulations - 2 by Siebenmann

3-Dimensions + knots

Group actions

Embedding & Unknotting in Higher Dimensions

Surgery

Participants in 1961

R. D. Anderson	G. R. Livesey
* J. Andrews	L. F. McAuley Jr.
B. J. Bell	E. E. Moise
R. H. Bing	R. H. Rosen
T. R. Brahana	D. E. Sanderson
* M. Brown	John Stallings
S. S. Cairns	G. T. Whyburn
Ed Connolly	Lucille Whyburn
M. L. Curtis	R. L. Wilder
P. Doyle	E. C. Zeeman
C. H. Edward Jr.	
* D. B. A. Epstein	* Hempel
M. K. Fort Jr.	Cantrell
R. H. Fox	
* Herman Gluck	
O. G. Harrold Jr.	
T. Homma	
J. Kister	
S. Kinoshita	

* Not retired

1977 1-12 August "Geometric Topology"
Ed Cantrell
Low dimensional manifolds
Topology of manifolds - Sullivan - Lipschitz
Shape theory & ∞ -dim topology - Chapman

1985 5-16 Aug "Geometry & Topology"
Ed McCrory & Shifrin

Some gauge theory, but not too
much low-dim topology due to
MSRI year in 84-85.

1993 2-13 Aug "Geometric Topology"
Ed W. Kazhdan
Mostly low dim top.

$$\tau_4(M) = \sum_{\theta} \rho^{\mu(M, \theta)}$$

where $\rho = e^{2\pi i(-2/16)}$, Pauli
Melvin

$\theta =$ spin structure on 3-manifold M

$\mu =$ μ -invariant $\in \mathbb{Z}/16$ "closed, oriented"

Let (M, θ) be spin boundary of W^4 .

Then $\mu(M, \theta) = \sigma(W^4) = \text{signature } W^4 / 16$

This is independent of choice of W because

$$\sigma(X^4) \equiv 0 \pmod{16}$$

if X^4 is smooth, closed, spin.
(Rohlin)

A spin structure θ is a trivialization of T_M over the 1-skeleton which extends over the 2-skeleton.

$$\{\theta\} \xrightarrow{1-1} H^2(M; \mathbb{Z}/2) \cong H_1(M; \mathbb{Z}/2)$$

∃? polynomial time algorithm to evaluate τ_4 ? (M Freedman)

Probably not!

The number of spin structures on M ,
" 2^{β_1} ", grows exponentially in $\beta_1 = \text{rank } H_1$
 $|H_1(M; \mathbb{Z}/2)|$

Don't want to find different spin W^4
for each θ on M^3 ,
so fix W^4 + then

$$\mu(M, \theta) = \sigma(W) - F \cdot F + 8 \text{Arf}(F) \quad (16)$$

where F^2 is a smooth embedded, orient
surface in W^4 such that θ extends
across $W - F$ but not across any
component of F .

F is P.dual to w_2

Arf(F) is the Arf invariant of a quadratic enhancement on $H_1(F; \mathbb{Z}/2)$

(If F non-orientable, then

$$\mu(M, \theta) = \sigma(W) - F \cdot F + 2 \text{Brown}(F) \quad (16)$$

How to find F? How to see θ ?

Let W^n be determined by a framed link L

Spin structures θ on $M = \partial W$ ^{n-components} correspond to "characteristic" sublinks C which satisfy $C \cdot L_i \equiv L_i \cdot L_i \pmod{2}$ (2) $\forall L_i \in L$
 linking framing

Construct F from cores of 2-handles attached to components of C, together with an oriented Seifert surface for C.

$$\text{Now } \tau_4(M) = \sum_{\Theta, C} \rho^{\sigma(w) - F \cdot F + 8 \text{Arf}(F)}$$

$\sigma(w)$ is a one-time calculation

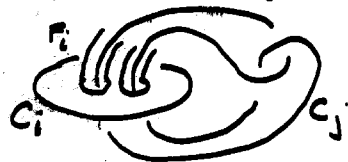
$F \cdot F$ is "quadratic", does not depend on F but only on linking form, & is not a problem.

But $\text{Arf}(F) \stackrel{\text{Arf}(C)}{\text{Arf}(F)}$ does depend on topology of F & we don't want to construct an F for each $\Theta = \text{char. link } C$.

Need "local" definition of $\text{Arf}(F)$.

Given C , suppose $\text{lk}(C_i, C_j) = 0 \forall C_i, C_j$

Then each C_i bounds a surface F_i ^{C_j} which is disjoint from the other C_j .



$$\text{Then } \text{Arf}(L) \equiv \sum_i \text{Arf}(C_i) - \sum_{i,j} h(C_i, C_j) + \sum_{i,j,k} T(C_i, C_j, C_k)$$

$$\text{Brown}(M_L) \equiv \sum_i \text{Brown}(C_i) - 2 \sum_{i,j} h(C_i, C_j) + 4 \sum_{i,j,k} T(C_i, C_j, C_k)$$

where: $T(C_i, C_j, C_k) =$ number, mod 2, of triple points in $F_i \cap F_j \cap F_k$.

$h(C_i, C_j) =$ framing, mod 2, of $F_i \cap F_j$ given by normals (to $F_i \cap F_j$) into F_i or F_j .

So, computing τ_4 in polyn time

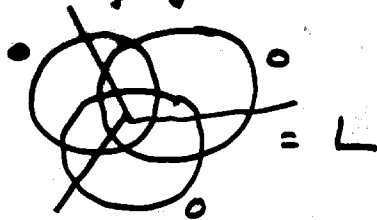
comes down to evaluating

$$\sum_C \rho \sum_{i,j,k} T(C_i, C_j, C_k) = \sum_C (-1)^{\rho(C)}$$

since $\text{Arf } C_i$ is a one time calculation for each component of L , + h is "quadratic".

0-surgery on Borromean rings gives T^3

T^3 has
8 spin structures



For $\gamma, \mu = 0$

For spin structure $C = L,$

$$\mu = 8 = 8 \text{Arf}(C) = 8T(L_1, L_2, L_3)$$

$$\text{so } \tau_4 = 7 - 1 = 6.$$

Let $g = \sum_{\substack{C, J, K \\ C, J, K}} a_{Cjk} x_C \cdot x_J \cdot x_K$ be a cubic form over $\mathbb{Z}/2$ with n variables $a_{Cjk} \in \mathbb{Z}/2$

Let L_g be the n strand unbraided but with the Borromean rings $(\sigma_i \sigma_j^{-1})^2$ "tied" in the braid in the $C, J, + K$ strands for each $a_{Cjk} = 1$.

$$\text{Then } \tau_4(M_g) = \sum_C p^{8A} = \sum_C (-1)^A \text{ for } A = \sum_{\substack{C, C, C_k \in C}} a_{Cjk}$$

$$\sum_{\substack{(x_1, \dots, x_n) \\ \in (\mathbb{Z}/2)^n}} (-1)^{g(x_1, \dots, x_n)} \quad (x_1, \dots, x_n) \in (\mathbb{Z}/2)^n$$

Evaluating $\sum_{(x_1, \dots, x_n)} (-1)^{\sum x_i}$ is #P-complete

P A problem is in the class P if \exists a polyn time algorithm giving a solution

NP A problem is NP (non-deterministic polyn time) if, e.g. a conjectured solution can be checked in polyn. time.

#P An NP problem is in #P if \exists polyn time algorithm to determine the number of solutions.

3-SAT is a universal #P problem

To count solutions to expressions

$E = A_1 \wedge A_2 \wedge \dots \wedge A_m$ where each A_k is of the form $(x_{l_1}^{z_1} \vee x_{l_2}^{z_2} \vee x_{l_3}^{z_3})$.

Each x_{l_i} is assigned T or F & want E to be True.

$$3\text{-SAT} \equiv \sum (-1)^{\sum (x_i - \bar{x}_i)}$$

change E to system of equations over \mathbb{Z}_2
by letting $T=1$, $F=0$, e.g.

$$(x_1 \vee x_2 \vee \bar{x}_3) \rightarrow (1-x_1)(1-x_2)x_3 = 0$$

change these cubic equations (one for each A_i)
to quadratic equations by
letting a new variable $x_{ij} = x_i \cdot x_j$

Thus $(1-x_1)(1-x_2)x_3 = 0$ becomes

$$x_{12} - x_1 \cdot x_2 = 0 \quad \& \quad (1-x_1-x_2+x_{12})x_3 = 0$$

So 3-SAT is equivalent to counting
the number of solutions over \mathbb{Z}_2 , $\#X$,
to k quadratic equations in m variables

$$X = \begin{cases} g_1(x_1, \dots, x_m) = 0 \pmod{2} \\ \vdots \\ g_k(x_1, \dots, x_m) = 0 \pmod{2} \end{cases}$$